

# Name/Identifiers and their 6 Attributes

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## → 1. Name

= identifier

length, legal chars, case-sensitivity, special words

can be one-one, many-one, or none-one mapping to memory

## → 2. Address

aliasing

point to a location in memory

may vary dynamically

anonymous - new (p)  
p<sup>^</sup> points to location that  
has no associated name

## → 3. Type

range of values + legal operations

variable, constant, label, pointer, program, ...

## → 4. Representation/Value

interpreted contents of the location

l-value (address)

r-value (value)

## → 5. Scope

Range of statements over which the variable is visible.

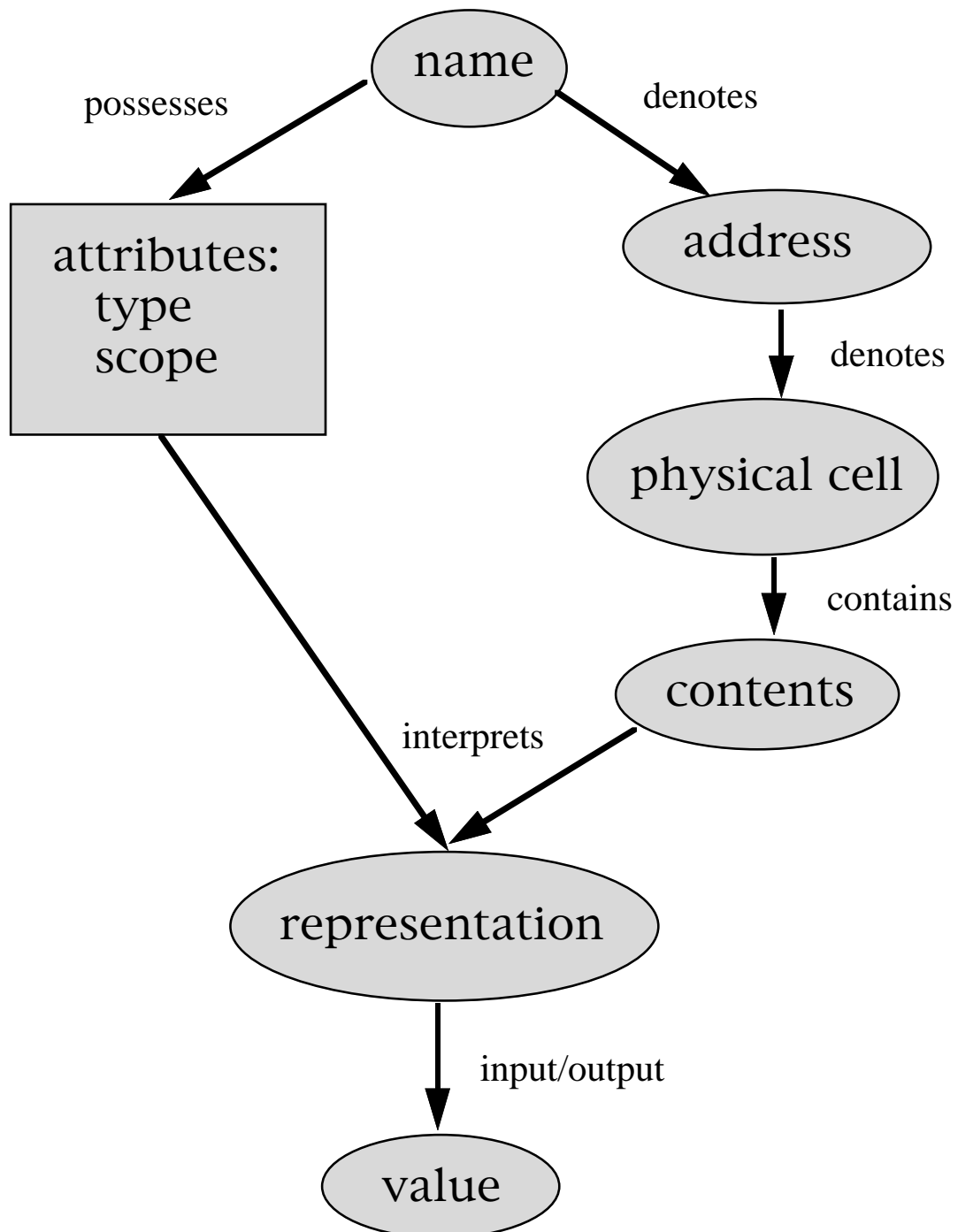
Static/dynamic

## → 6. Lifetime

Time during which the variable is bound to a storage location.

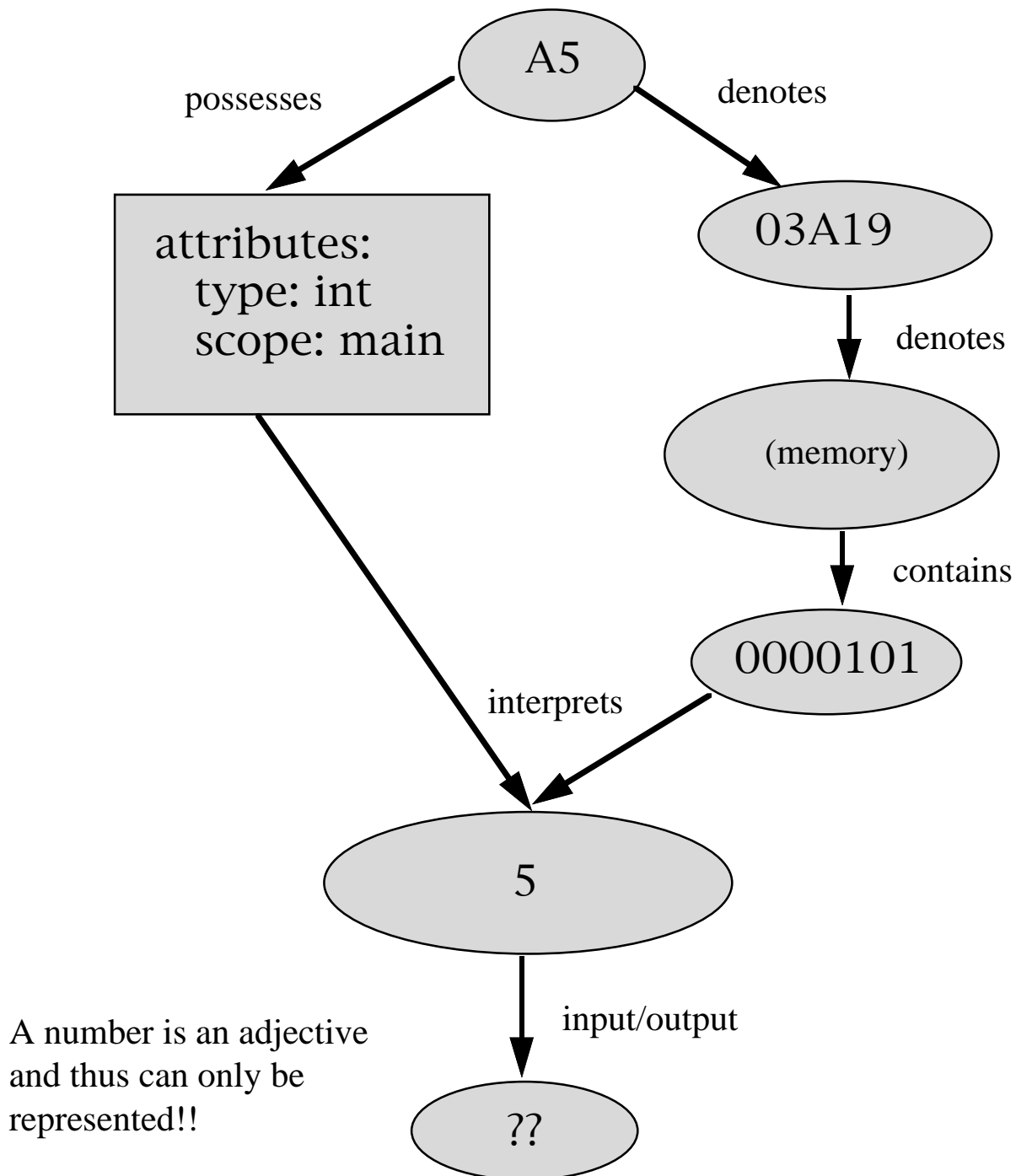
# Name/Value Graph

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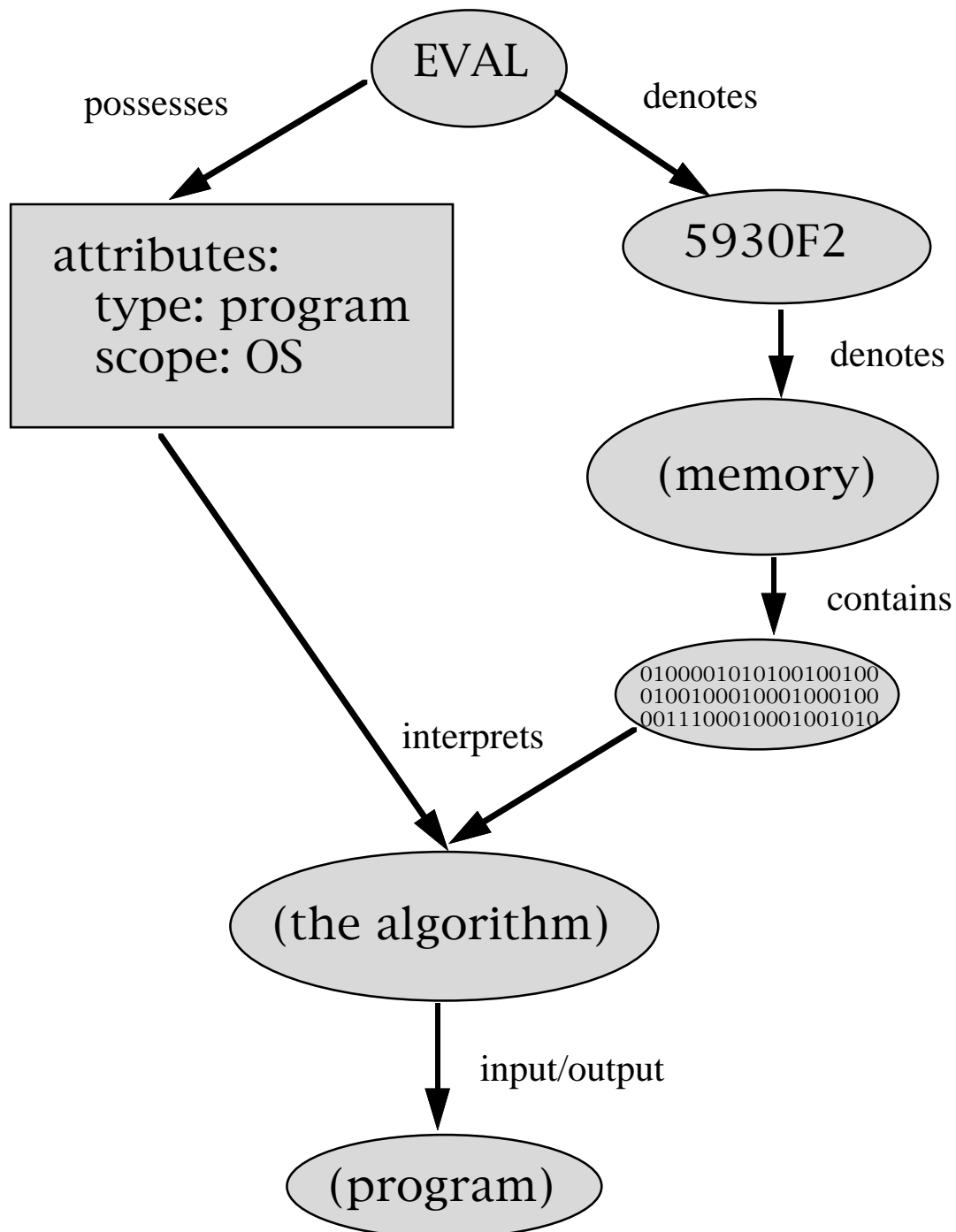
# Name/Value Graph - Example I

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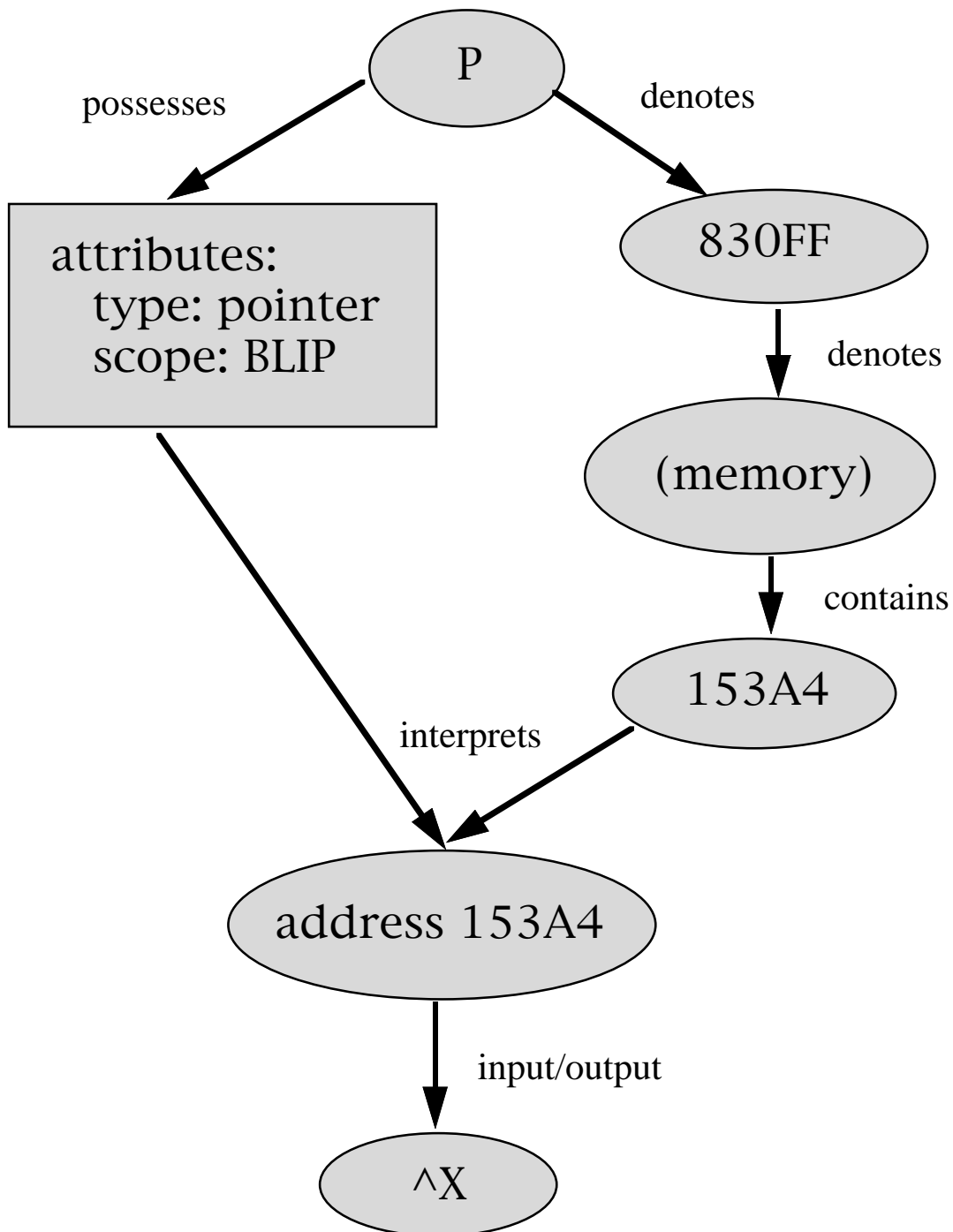
## Name/Value Graph - Example II

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## Name/Value Graph - Example III

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# Binding

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- **How and when are attributes bound to identifiers?**
  - ***Static***
    - occurs before runtime (compile time, link/load time)
    - constant throughout program execution
  - ***Dynamic***
    - occurs or can change during runtime
- **In many ways, the various binding times determine the flavor of a language.**
- **As binding time gets earlier:**
  - **efficiency goes up**
  - **safety goes up**
  - **flexibility goes down**

# BINDING TIMES

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- LANGUAGE DESIGN
- COMPILER DESIGN/IMPLEMENTATION
- COMPILE TIME
- LINK/LOAD TIME
- RUN TIME

# BINDING TIMES

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|                                     | Language Design | Compiler Design | Compiler Time | Link/Load Time | Run Time |
|-------------------------------------|-----------------|-----------------|---------------|----------------|----------|
| Type                                |                 |                 |               |                |          |
| Set of objects in Type              |                 |                 |               |                |          |
| Association of Type with Identifier |                 |                 |               |                |          |
| Initial value “of” identifier       |                 |                 |               |                |          |
| Possible Meanings of Operators      |                 |                 |               |                |          |
| Representation of Values            |                 |                 |               |                |          |



# Type Binding

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**When is type bound to variable?**

**How is binding specified?**

- ***Static typing* (before runtime)**

- **explicit declaration**

- var x: integer

- **implicit declaration by prior agreement**

- e.g., Fortran:

- I = 5      First mention of I -- integer

- PI = 3.14    First mention of PI -- real

- **Advantages:**

- **cheaper**

- **safer**

- **Disadvantage:**

- **less flexible**

# Type Binding

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- ***Dynamic Binding*** (after compile time)
  - Identifier gets type of value assigned to it as needed.  
    `x := 5 -- x is of type integer`  
    ...  
    `x := "foo" -- x is of type string`
- **Advantage:**
  - flexibility
- **Disadvantages:**
  - runtime overhead -- extra work at run-time
  - poor error detection -- type checking is replaced by type changing
  - you will lose 5 points in an homework assignment for reusing identifiers
- **More about types later. . .**

# Scope

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- ***Static (lexical) scope***

- **Scope of an identifier is determined by the textual layout of the program.**
- **In block structured languages, scope of an identifier is**
  - the unit in which it is defined, plus
  - all units immediately nested inside the declaring unit (excluding those in which the variable is redeclared), plus
  - all units within which the declaring unit is nested.
- **To find the declaration of an identifier,**
  - (1) lift all declarations to the top of the unit
  - (2) look through the statically enclosing units until a declaration is found.

- ***Dynamic Scope***

- **Scope of an identifier depends on program execution, and therefore changes dynamically.**
- **To find declaration, look up through the call chain.**

## Example (evaluate both ways)

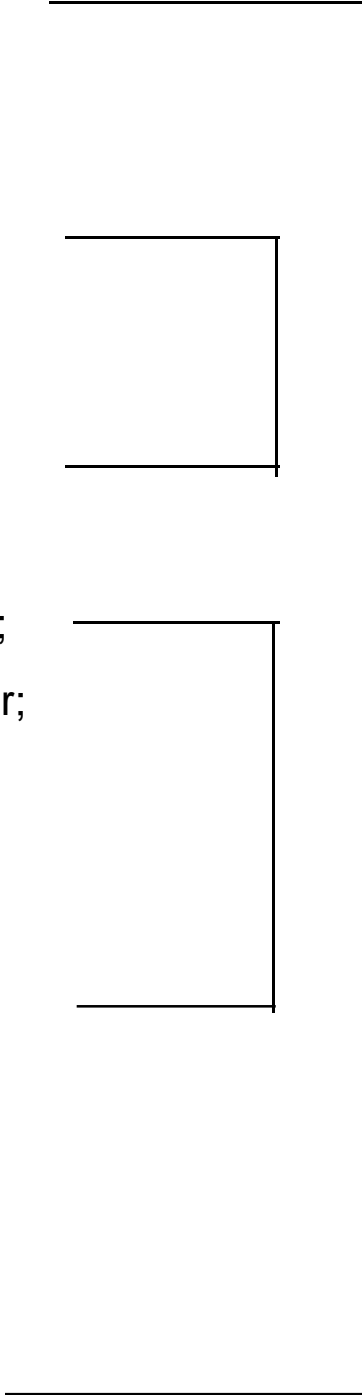
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```
program foo;
var x: integer;

    procedure f;
    begin
        print(x);
    end f;

    procedure g;
    var x: integer;
    begin
        x := 2;
        f;
    end g;

begin
    x := 1;
    g;
end foo.
```



# Lifetime (= extent)

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- **The lifetime of a variable is the interval of time during which it is bound to a specific memory location.**
- **Static identifiers**
  - bound to memory cells before execution (load time)
  - retain same binding throughout execution
  - efficient, inflexible
  - allow history-sensitivity
  - do not support recursion
  - Ex: FORTRAN variables
- **Semidynamic identifiers**
  - storage allocated when unit is called
  - storage deallocated when unit returns
  - allows recursion
  - Ex: Variables declared in Pascal Procedures

# Lifetime (continued)

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- **Explicit Dynamic Identifiers**

- storage allocated and deallocated by programmer
- *new, dispose* in Pascal
- flexible and efficient, but dangerous
- Ex: Pointer variables

- **Implicit Dynamic Identifiers**

- automatically bound to storage as needed
- storage automatically reclaimed when no longer needed
- flexible, safe, less efficient
- Ex: lists in LISP, Prolog

# Scope $\neq$ Lifetime

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- **lifetime > scope: storage that can't be accessed through that variable.**

```
var p: ^integer;  
begin  
    ...  
    new(p)  
    ...  
end
```

└─ here, storage is still allocated but p is not defined,  
 └─ Lifetime > Scope

- **scope > lifetime: variable without storage.**

```
var p: ^integer;  
begin  
    ...  
    new (p)                addr   i   "name"  
                           of mem  
    ...  
    dispose(p)  
    └─ here, p is defined but has no value,  
       └─ Scope > Lifetime  
end
```

## Scope $\neq$ Lifetime (continued)

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- **Also, scope has "holes" during execution, but lifetime does not.**

```
procedure f;
```

```
var x: real;
```

```
begin
```

```
    ...
```

```
end
```

```
procedure g;
```

```
var x: integer;
```

```
begin
```

```
    ...
```

```
    f
```

```
    ...
```

```
end
```

out of x's scope during  
execution of f (assuming static  
scope), but x's lifetime persists.



# Variable Initialization

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- **static or dynamic**
- **once for static variables, at each allocation for dynamic variables**

- **many possible methods:**

sum: integer := 0;      (Ada)

int first := 10      (Algol 68 initialization)

int first = 10      (Algol 68 constant declaration)

- **unavailable in Pascal**
- **default initializations**

# Aliasing

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- **Two variables are aliases if they share the same storage.**

|                    |                   |
|--------------------|-------------------|
| var x,y: ^integer; | var x,y: integer; |
| begin              | begin             |
| new (x);           | x := 5;           |
| x^ := 5;           | y := x;           |
| y := x;            | x := 10;          |
| x^ := 10;          | writeln (y);      |
| writeln (y^);      | end               |
| end                |                   |

- **Also results from *var* parameters:**

```
procedure p (var x,y: integer);  
....  
p (a,b);  
p (x1,x1);
```