

Name/Identifiers and their 6 Attributes

→ 1. Name

= identifier

length, legal chars, case-sensitivity, special words

can be one-one, many-one, or none-one mapping to memory

→ 2. Address

aliasing

point to a location in memory

may vary dynamically

anonymous - new (p)
p[^] points to location that
has no associated name

→ 3. Type

range of values + legal operations

variable, constant, label, pointer, program, ...

→ 4. Representation/Value

interpreted contents of the location

l-value (address)

r-value (value)

→ 5. Scope

Range of statements over which the variable is visible.

Static/dynamic

→ 6. Lifetime

Time during which the variable is bound to a storage location.

Binding

- **How and when are attributes bound to identifiers?**
 - ***Static***
 - occurs before runtime (compile time, link/load time)
 - constant throughout program execution
 - ***Dynamic***
 - occurs or can change during runtime
- **In many ways, the various binding times determine the flavor of a language.**
- **As binding time gets earlier:**
 - **efficiency goes up**
 - **safety goes up**
 - **flexibility goes down**

BINDING TIMES

- LANGUAGE DESIGN
- COMPILER DESIGN/IMPLEMENTATION
- COMPILE TIME
- LINK/LOAD TIME
- RUN TIME

BINDING TIMES

	Language Design	Compiler Design	Compiler Time	Link/Load Time	Run Time
Type					
Set of objects in Type					
Association of Type with Identifier					
Initial value “of” identifier					
Possible Meanings of Operators					
Representation of Values					

Type Binding

When is type bound to variable?

How is binding specified?

- ***Static typing* (before runtime)**

- **explicit declaration**

- var x: integer

- **implicit declaration by prior agreement**

- e.g., Fortran:

- I = 5 First mention of I -- integer

- PI = 3.14 First mention of PI -- real

- **Advantages:**

- **cheaper**

- **safer**

- **Disadvantage:**

- **less flexible**

Type Binding

- ***Dynamic Binding*** (after compile time)
 - Identifier gets type of value assigned to it as needed.
 `x := 5 -- x is of type integer`
 ...
 `x := "foo" -- x is of type string`
- **Advantage:**
 - flexibility
- **Disadvantages:**
 - runtime overhead -- extra work at run-time
 - poor error detection -- type checking is replaced by type changing
- **More about types later. . .**

Scope

- ***Static (lexical) scope***

- **Scope of an identifier is determined by the textual layout of the program.**
- **In block structured languages, scope of an identifier is**
 - the unit in which it is defined, plus
 - all units immediately nested inside the declaring unit (excluding those in which the variable is redeclared), plus
- **To find the declaration of an identifier,**
 - (1) lift all declarations to the top of the unit
 - (2) look through the statically enclosing units until a declaration is found.

- ***Dynamic Scope***

- **Scope of an identifier depends on program execution, and therefore changes dynamically.**
- **To find declaration, look up through the call chain.**

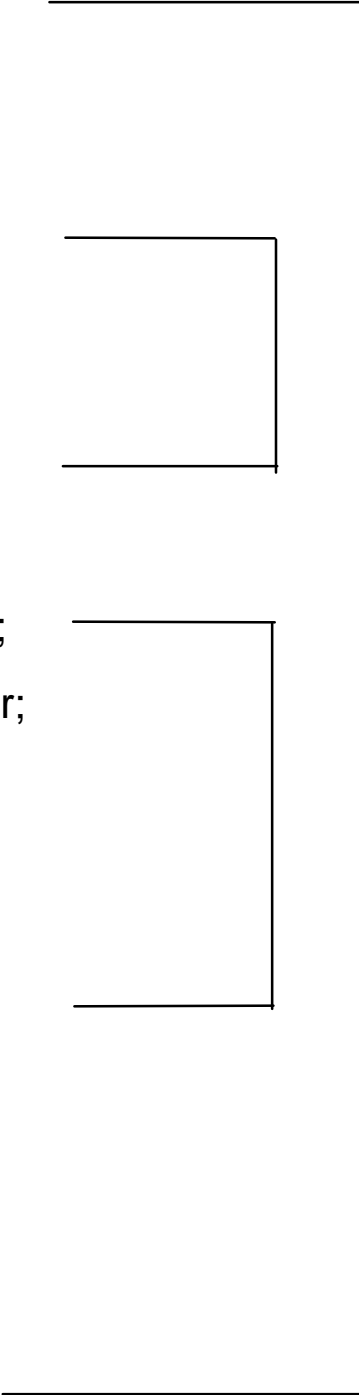
Example (evaluate both ways)

```
program foo;
var x: integer;

    procedure f;
    begin
        print(x);
    end f;

    procedure g;
    var x: integer;
    begin
        x := 2;
        f;
    end g;

begin
    x := 1;
    g;
end foo.
```



The diagram illustrates the nested scopes of the code. A large vertical line on the right side of the code blocks represents the scope of the program 'foo'. Inside this scope, there are two smaller vertical lines representing the scopes of procedures 'f' and 'g'. The scope of 'f' is nested within the scope of 'foo', and the scope of 'g' is also nested within the scope of 'foo'. The scope of 'g' is further nested within the scope of 'foo'.

Lifetime (= extent)

- The lifetime of a variable is the interval of time during which it is bound to a specific memory location.
- Static variables
 - bound to memory cells before execution (load time)
 - retain same binding throughout execution
 - efficient, inflexible
 - allow history-sensitivity
 - do not support recursion
 - Ex: FORTRAN variables
- Stack- dynamic variables
 - storage allocated when unit is called
 - storage deallocated when unit returns
 - allows recursion
 - Ex: Variables declared in Pascal Procedures

Lifetime (continued)

- **Explicit Heap - Dynamic Variables**
 - storage allocated and deallocated by programmer
 - *new, dispose* in Pascal
 - flexible and efficient, but dangerous
 - Ex: Pointer variables
- **Implicit Heap - Dynamic Variables**
 - automatically bound to storage as needed
 - storage automatically reclaimed when no longer needed
 - flexible, safe, less efficient
 - Ex: lists in LISP, Prolog

Scope \neq Lifetime

- **lifetime > scope: storage that can't be accessed through that variable.**

```
var p: ^integer;  
begin  
    ...  
    new(p)  
    ...  
end
```

↓ here, storage is still allocated but p is not defined,
Lifetime > Scope

- **scope > lifetime: variable without storage.**

```
var p: ^integer;  
begin  
    ...  
    new (p)                addr   i   "name"  
                           of mem  
    ...  
    dispose(p)  
    ↓ here, p is defined but has no value,  
      Scope > Lifetime  
end
```

Scope \neq Lifetime (continued)

- **Also, scope has "holes" during execution, but lifetime does not.**

```
procedure f;
```

```
begin
```

```
...
```

```
end
```

```
procedure g;
```

```
var x: integer;
```

```
begin
```

```
...
```

```
  f
```

```
...
```

```
end
```

out of x's scope during
execution of f (assuming static
scope), but x's lifetime persists.